



Micro-measurements in high capacity
networks

21. NORDUnet Network Conference

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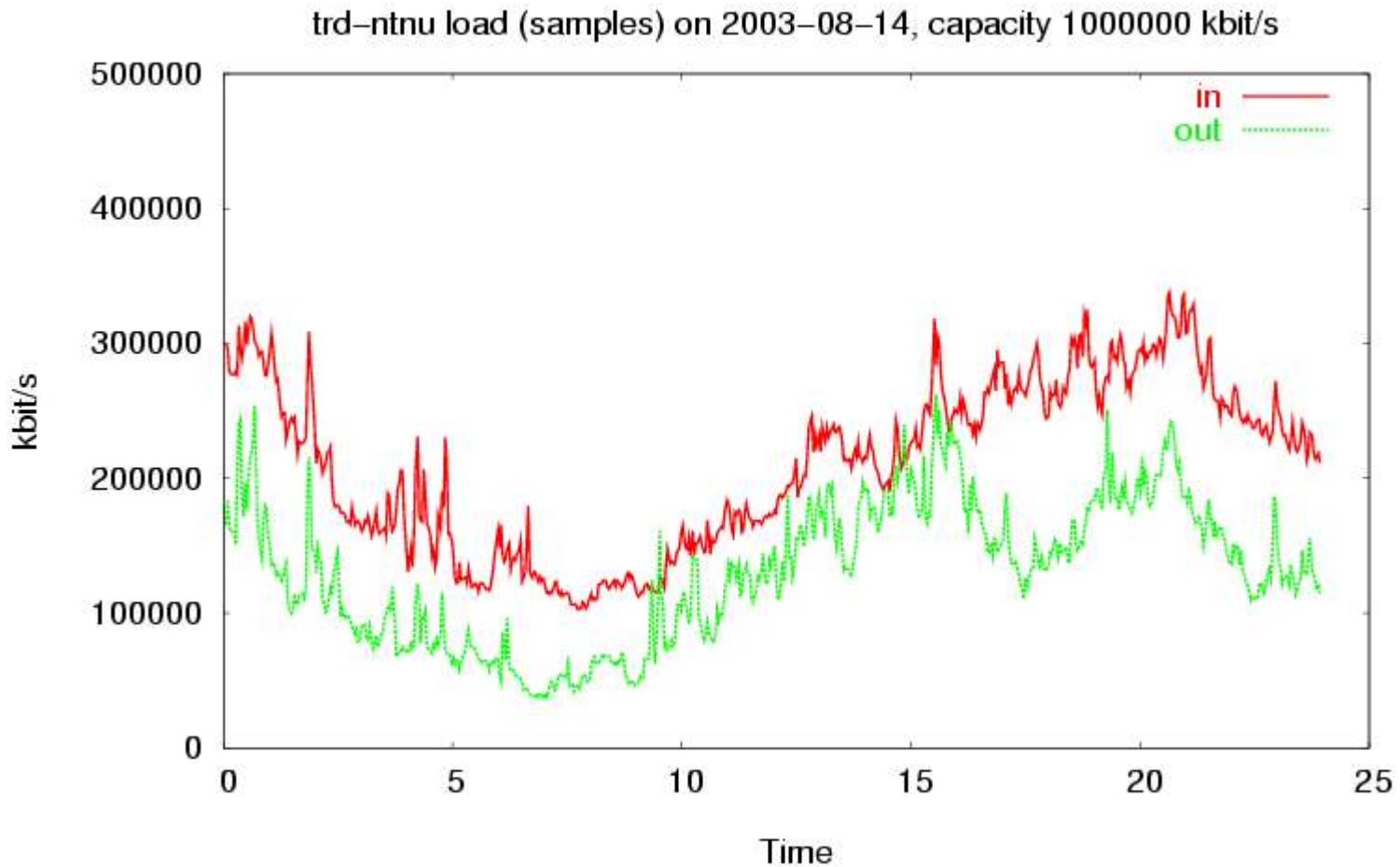
What do we know about our traffic?

- We monitor link traffic
- UNINETT statistics:
 - Year, month, week, day
 - Resolution: week – 2 minutes

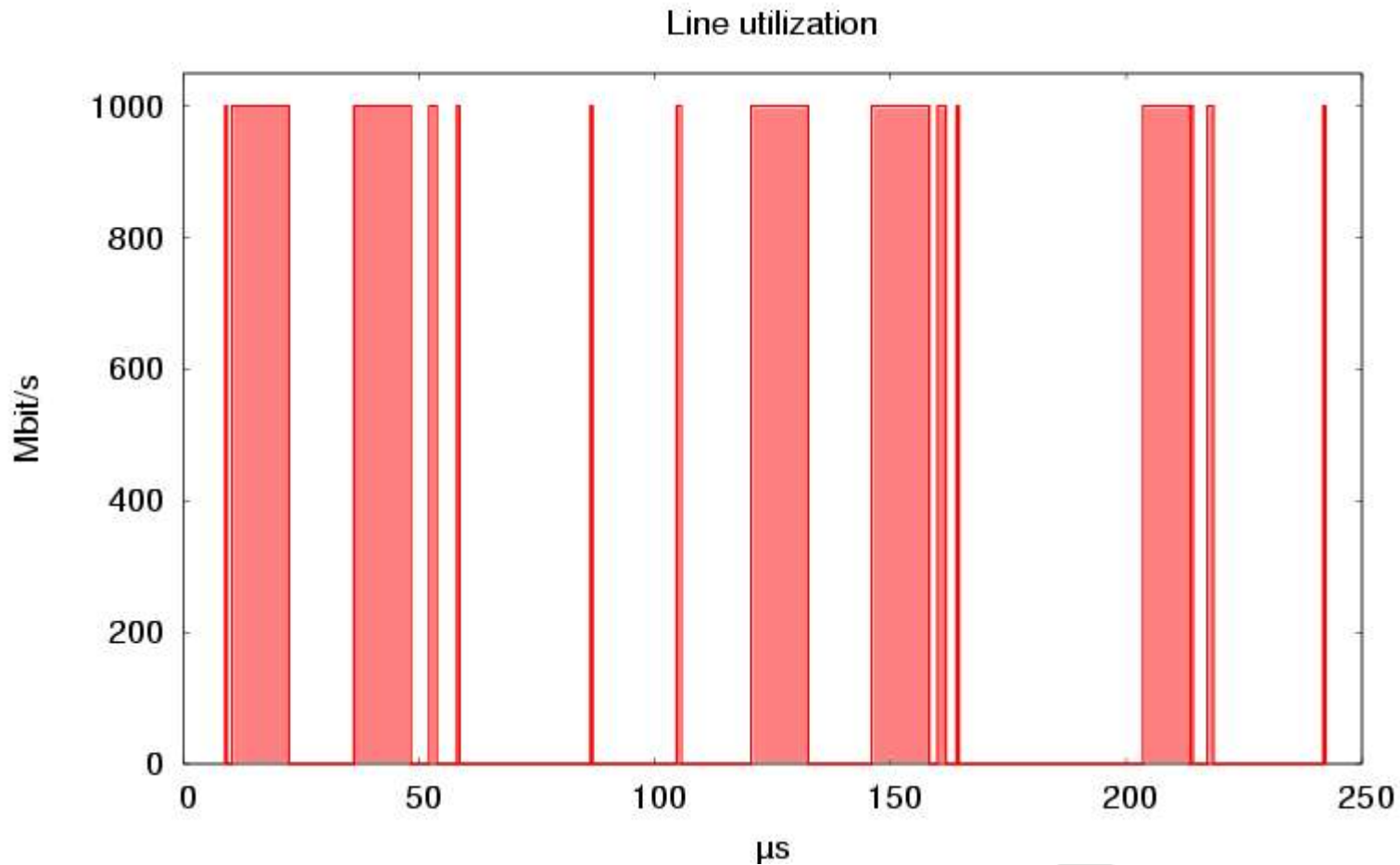
<http://drift.uninett.no/stats/tbl-report.html>

[http://drift.uninett.no/kartg/last/uninett/
norge/geo/nuh](http://drift.uninett.no/kartg/last/uninett/norge/geo/nuh)

A day in the life of a gigabit link



The microscopic view



The load is low. Are we safe?

- What happens between the samples?
- Network traffic is bursty
 - Models which assume statistical independence are too optimistic
- Self similarity has been observed
 - Traffic patterns are said to look the same when studied at very different time scales
 - Self similar traffic sources do not aggregate well
 - Studies from the 90's. Still relevant?

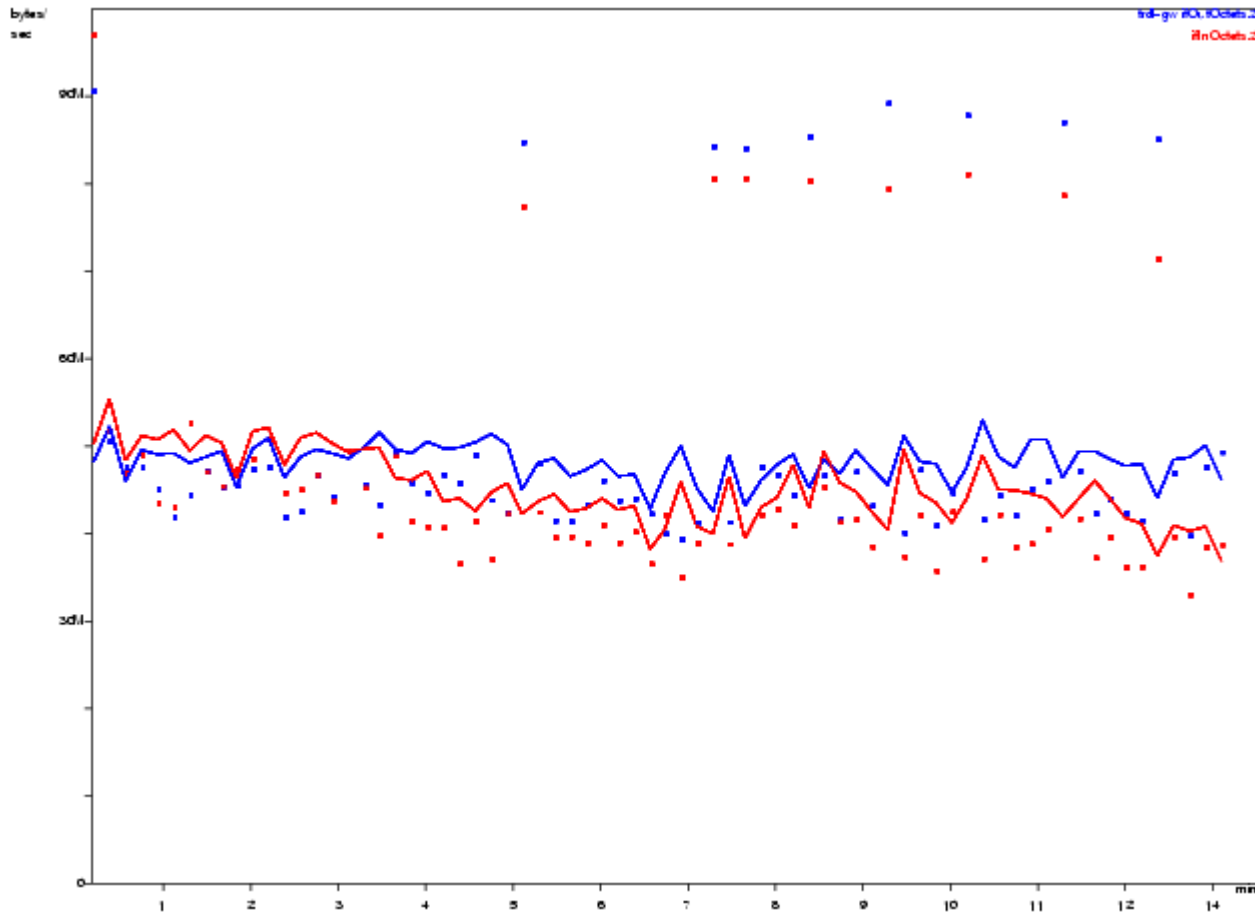
Measurement tools

- **SNMP counters**
Used to make UNINETT traffic statistics
- **Hardware probes**
 - **DAG cards from Endace**
Up to 2.5 Gbit/s
UNETT has 3
 - **COMBO-6, Masaryk University**
In development, part of SCAMPI project.

SNMP polling with a twist

- “Di-daaah” polling
 - For each regular sample, also take a shorter sample
- But
 - SNMP counters are updated infrequently
 - Updating them has low priority
- So on Cisco core routers, measuring $< 5s$ yields unreliable results

The cloud should be around the lines



What went wrong?

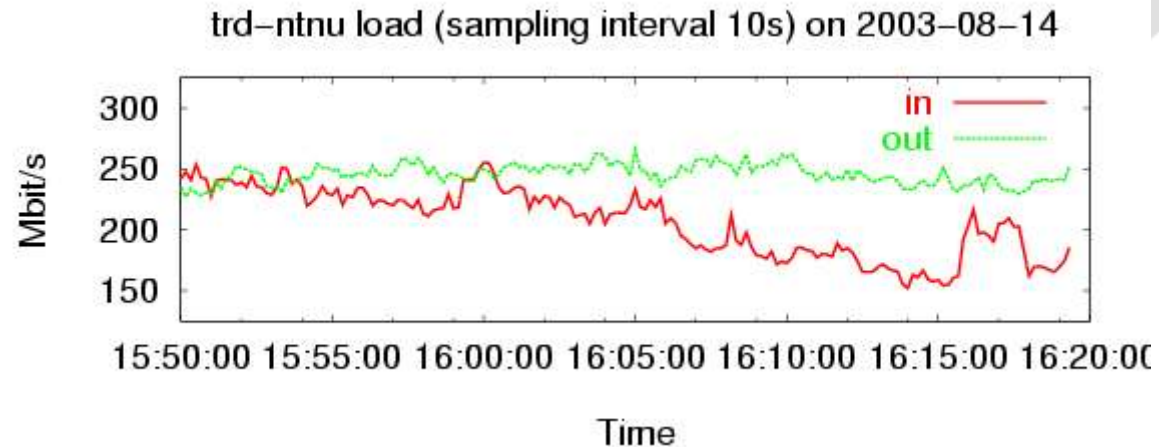
- Imagine that router updates counters once every long sampling period.
- Two cases
 - Update outside short sampling period
 - looks like zero traffic
 - Update inside short sampling period
 - looks like all traffic happens here
- Another idea: measure queues, but same problem

Is there a better way?

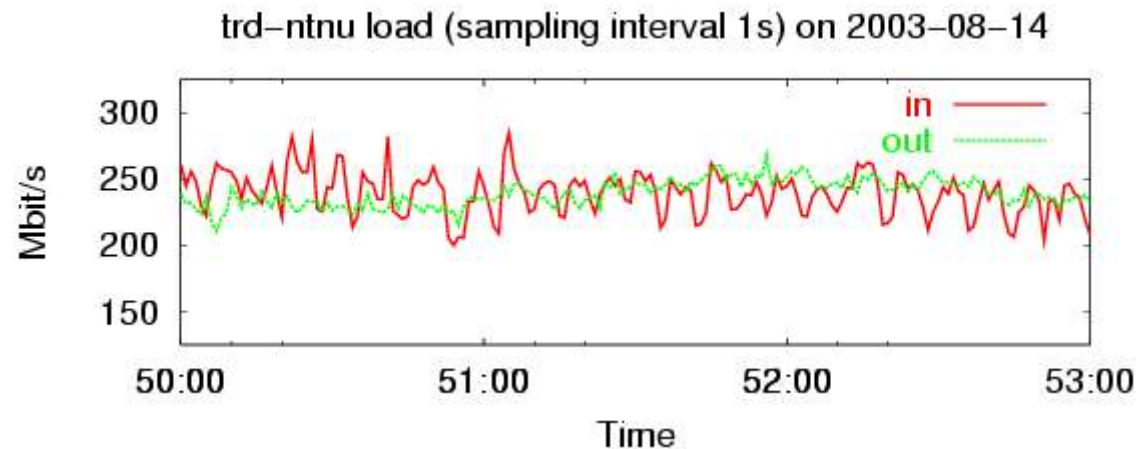
- Hardware probes
- PC card
 - Time source
 - Crystal
 - GPS
 - Optical splitter
- UNINETT uses DAG from Endace
- Masaryk University is developing Combo 6

Zooming in, shortening periods

10 s

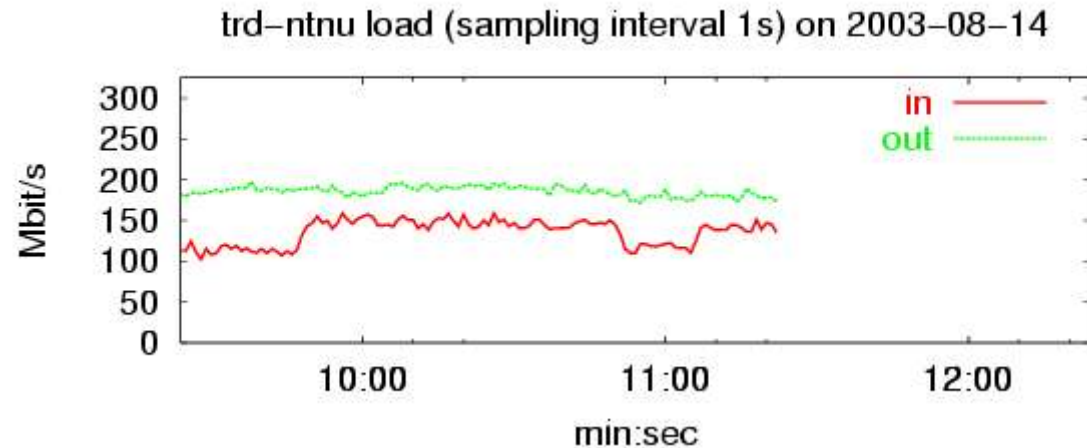


1 s

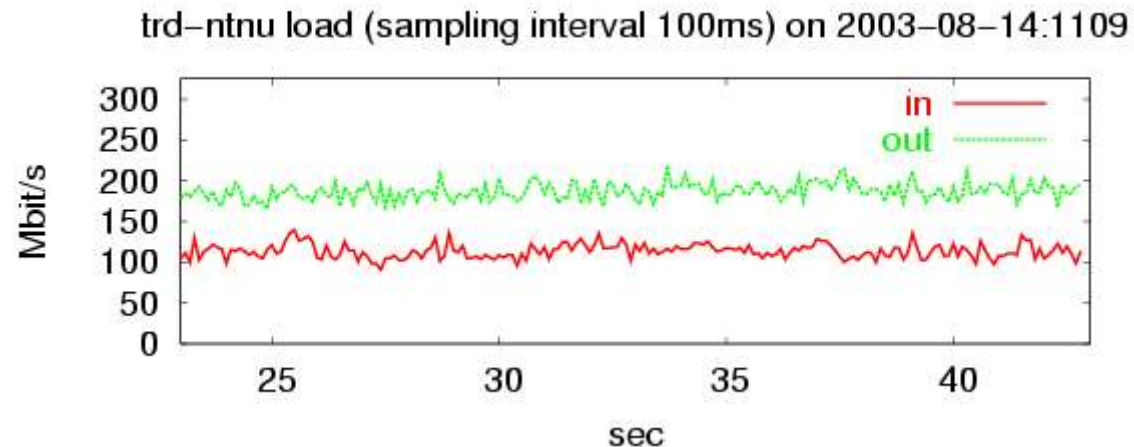


Zooming in (different trace)

1 s



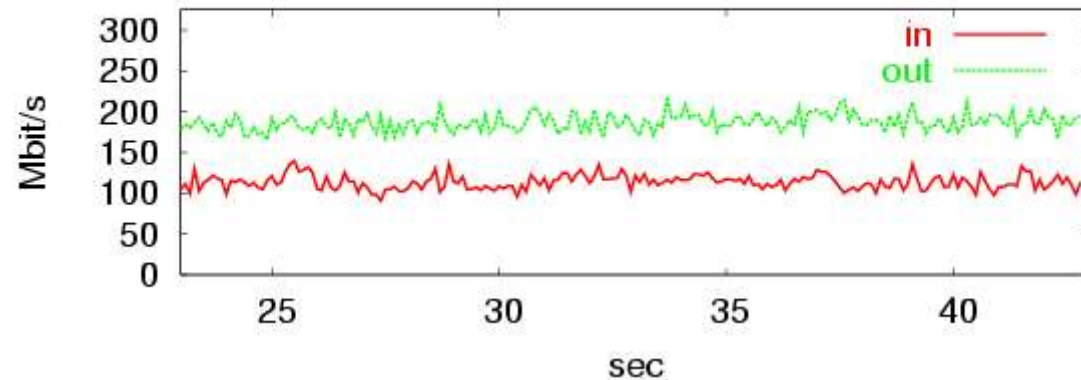
100 ms



Still zooming

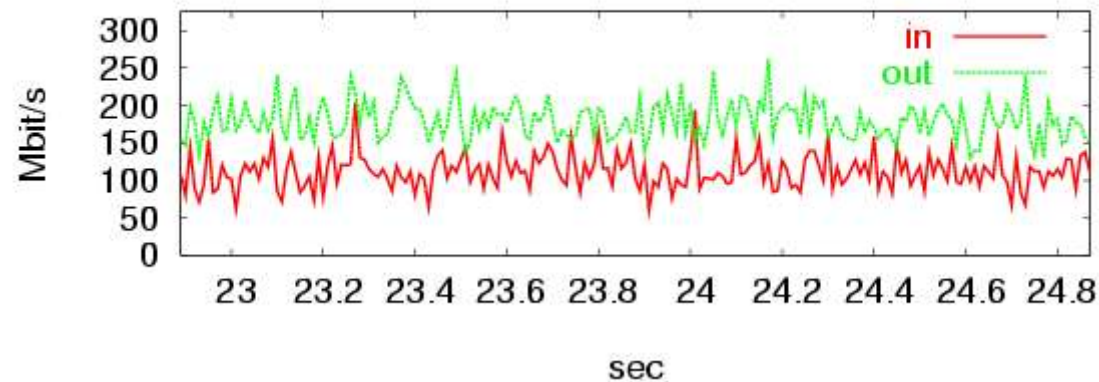
100 ms

trd-ntnu load (sampling interval 100ms) on 2003-08-14:1109



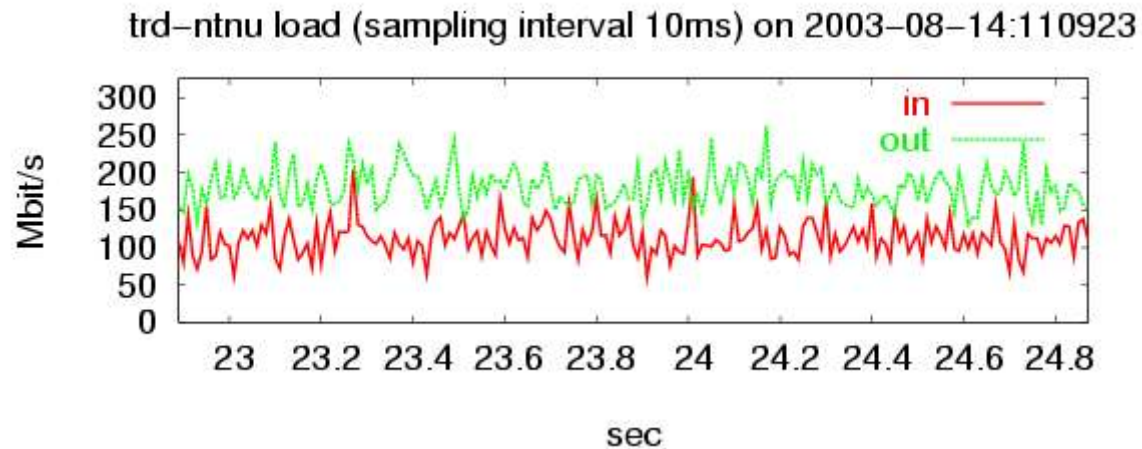
10 ms

trd-ntnu load (sampling interval 10ms) on 2003-08-14:110923

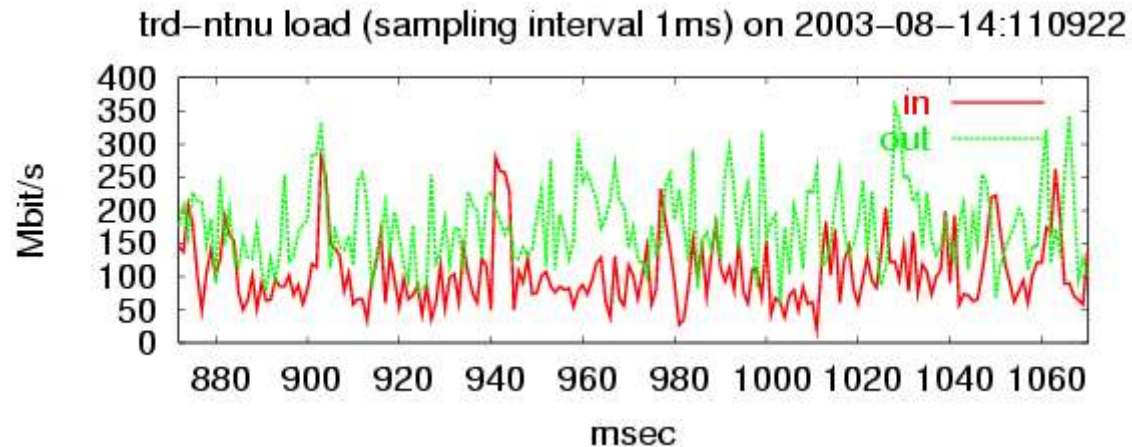


Are you bored yet?

10 ms

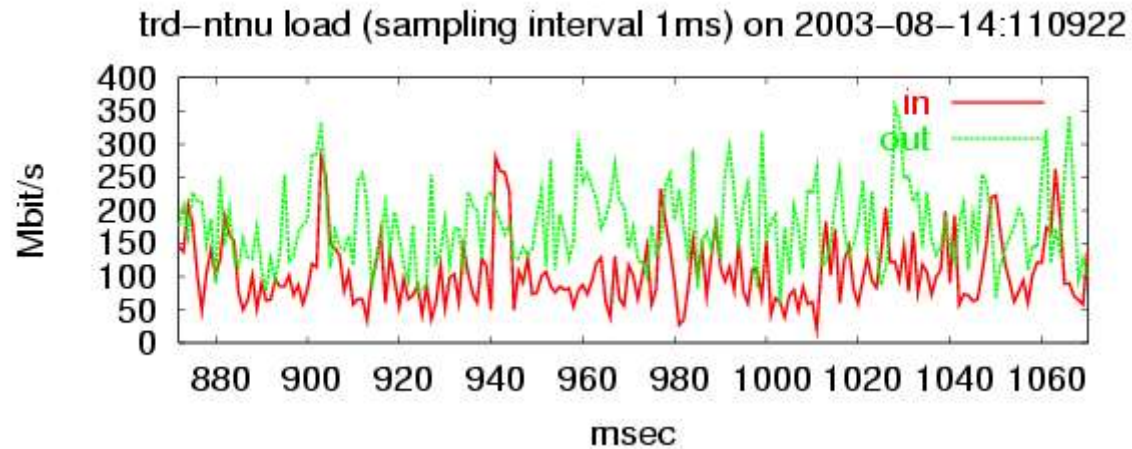


1 ms

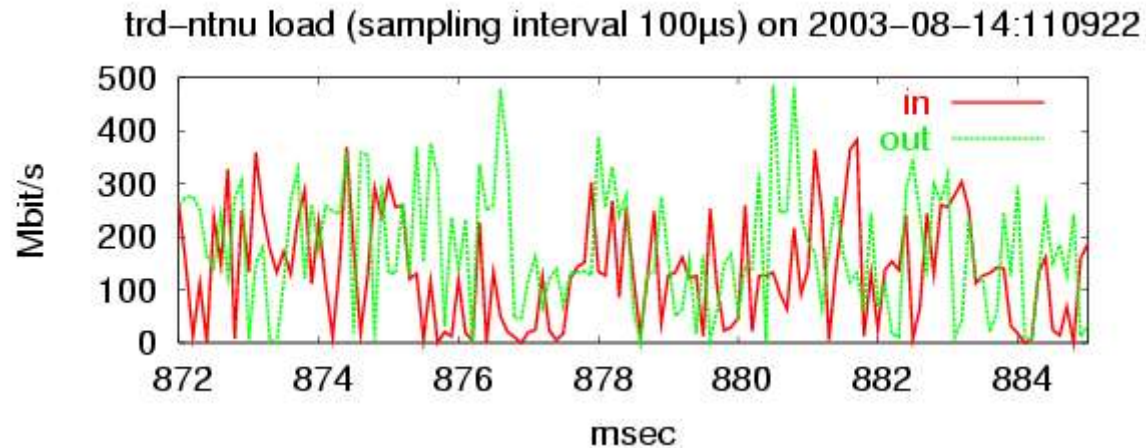


As short as it makes sense to go

1 ms



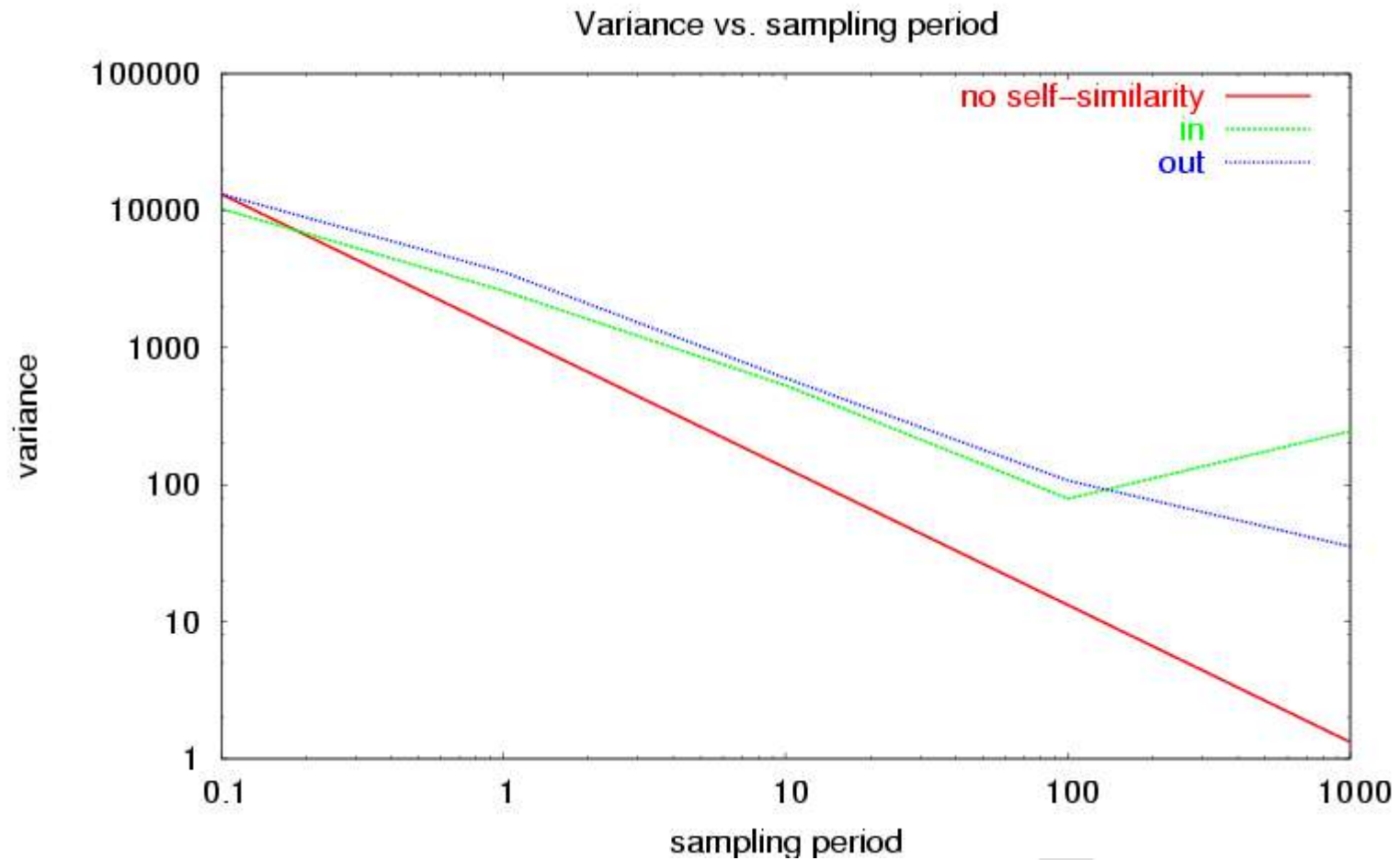
100 μ s



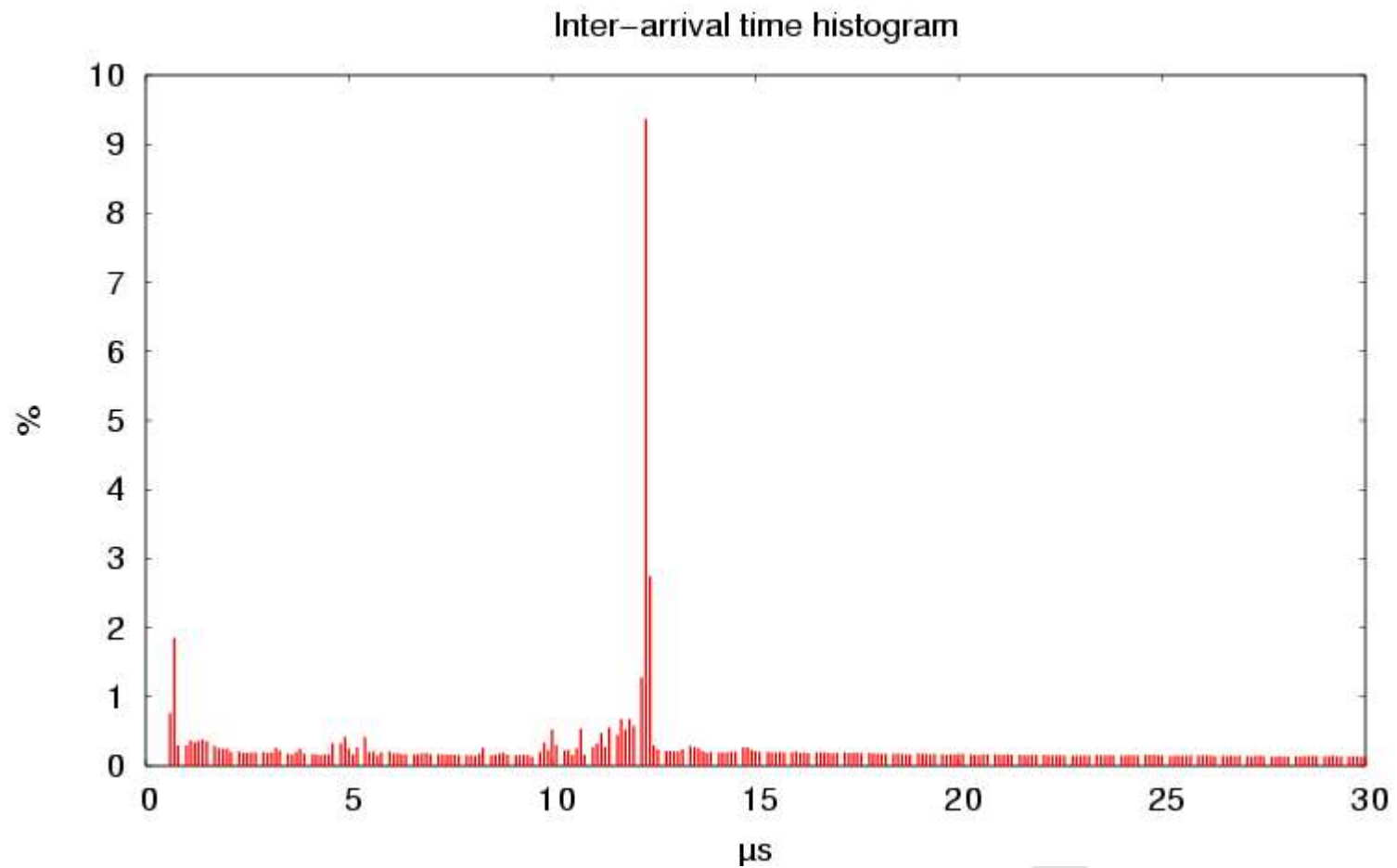
Is this traffic self-similar?

- That's how it looks
- But optical illusion is possible
 - Each plot shows the same number of samples
 - Consecutive samples are different
 - Upward and downward change equally likely
- However that may be, it is certainly bursty

It is indeed self-similar!



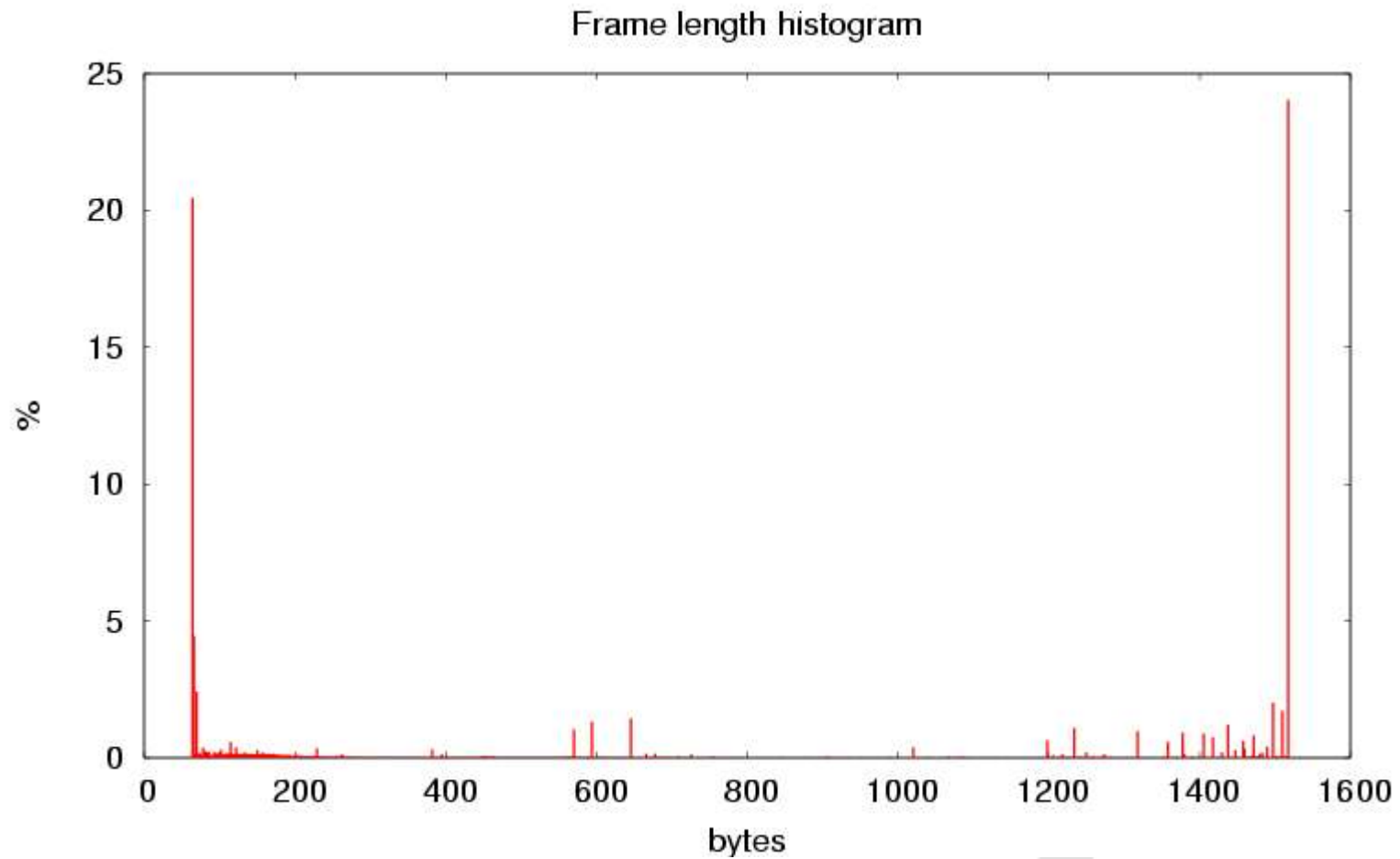
Inter-arrival times



Inter-arrival times

- Two peaks, max and min frame length
13.8% $\approx 12.3 \mu\text{s}$
13.0% 0.7-0.9 μs
- Min or max length frames sent just after another frame

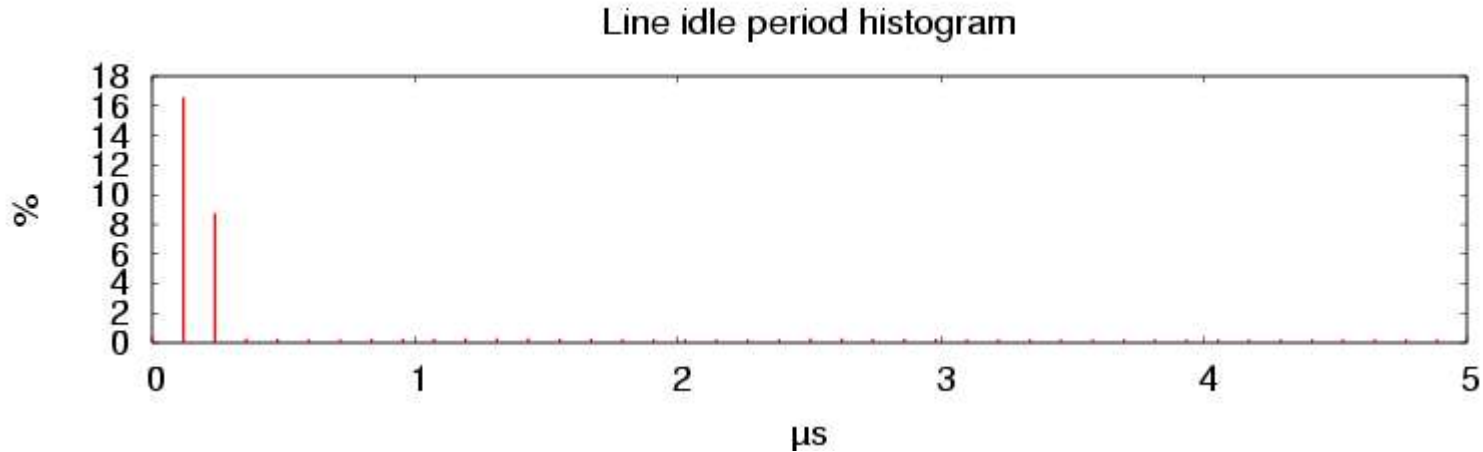
Frame lengths



Spikes

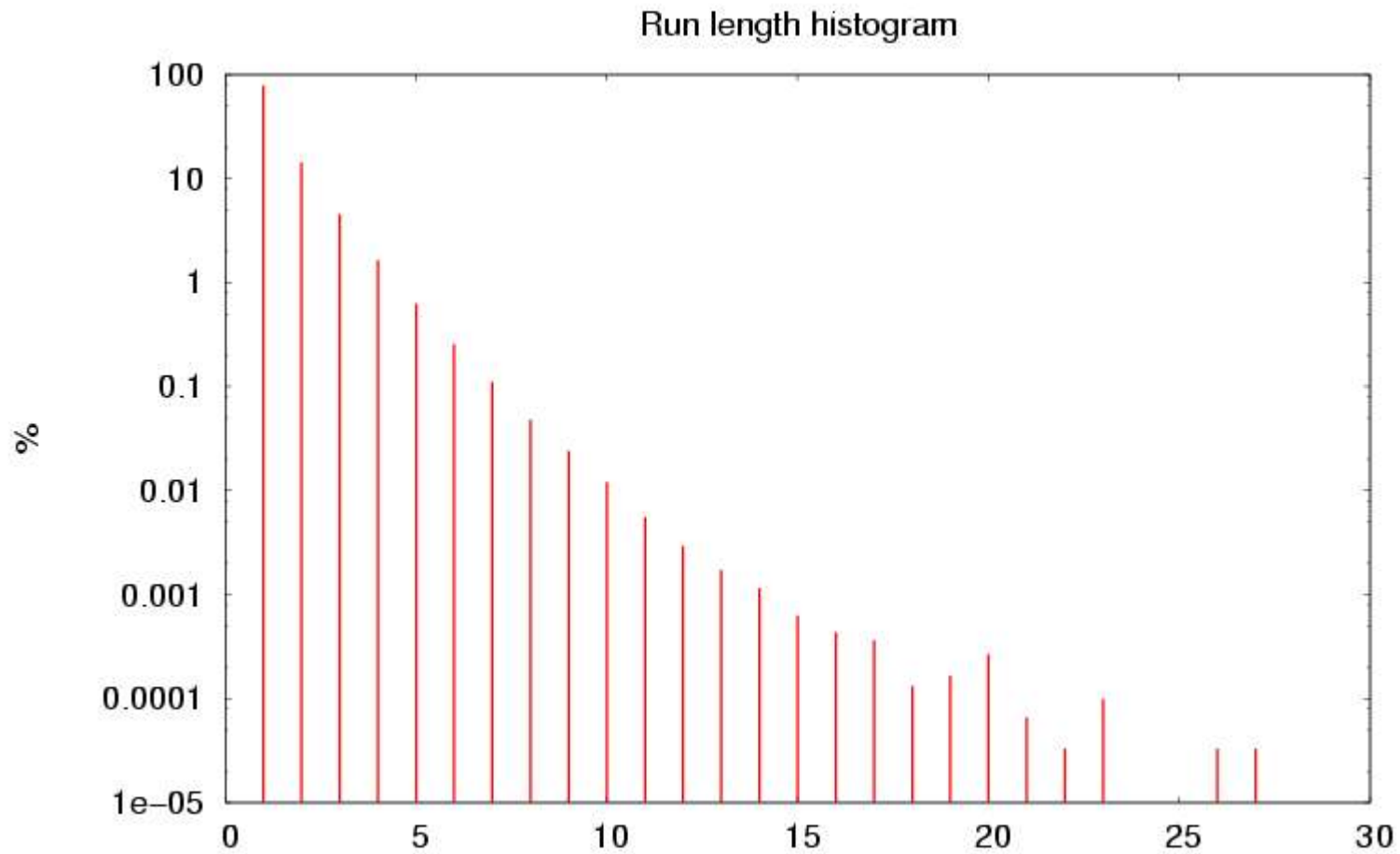
- 24% of frames are 1518 bytes
- 27% are 64 - 70 bytes
- Load average: 18.2%
- If arrivals were independent, spike for 12.3 μ s inter-arrival time would be $24\% \times 18\% = 4.3\%$
- Observed: 13.8%

Line idle time



- Very frequent values $< 0.3 \mu\text{s}$ must be line card sending frames as close to each other as it can.

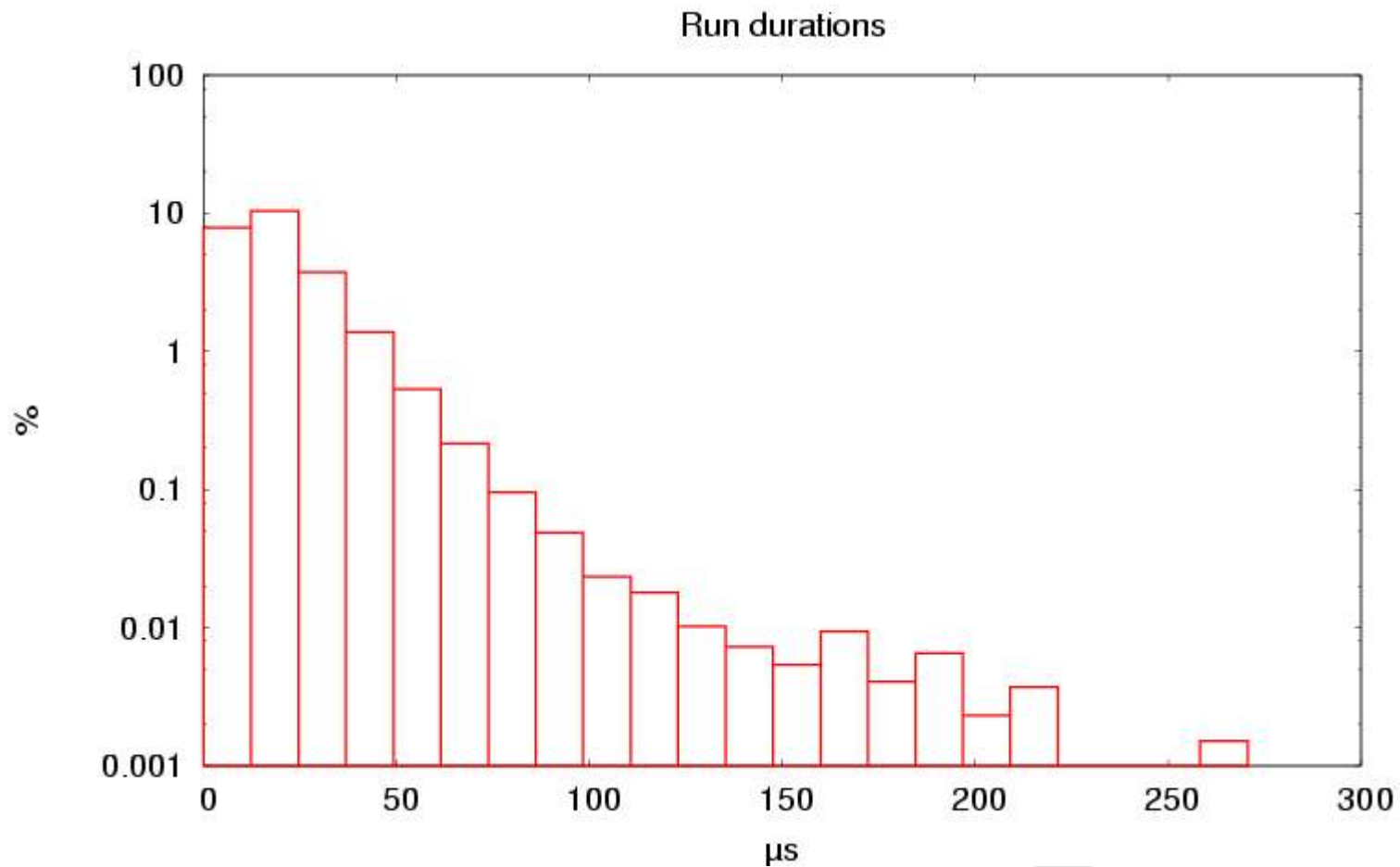
Runs of packets (log scale)



Runs of packets

- May indicate queuing
- 25% of frames are not first in a run
- Longest run in 2 minute period:
27 frames
34033 bytes
0.27 ms
- 99.99% of runs shorter than 10 frames

Run durations (log scale)



Conclusions

- Hardware probes reveal a lot
- Monitoring run lengths may be particularly useful.
 - Runs/queues may be long enough that we notice, but not yet so long that users notice

Thank you for your attention

- <http://domen.uninett.no/~jk/micromeasurements/>
- Email: Jon.Kare.Hellan@uninett.no